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## Lecture 8: <br> Item-to-item; Page Rank

## SI583: Recommender Systems

## Item-Item Collaborative Filtering

High-level approach:

- For each item $X$ find similar items $Y, Z$.
- For user Joe, recommend items most similar to items Joe has already liked


## Users-by-Items Matrix



## Normalize the Rows for User-User

## Algorithm

$$
\begin{aligned}
& X_{i J}=R_{i J}-R_{i} \\
& X=\left|\begin{array}{ccc}
1 / 3 & -2 / 3 & 1 / 3 \\
-2 / 3 & 1 / 3 & 1 / 3 \\
1 / 3 & -2 / 3 & 1 / 3 \\
-2 / 3 & 1 / 3 & 1 / 3
\end{array}\right| \quad R=\left|\begin{array}{lll}
1 & 0 & 1 \\
0 & 1 & 1 \\
1 & 0 & 1 \\
0 & 1 & 1
\end{array}\right|
\end{aligned}
$$

## Normalize the Columns for Item-Item

## Algorithm

$$
\begin{array}{cc}
W_{j k}=R_{j k} & -W_{k} \\
X=\left|\begin{array}{ccc}
0.5 & -.5 & 0 \\
-5 & 0.5 & 0 \\
0.5 & -.5 & 0 \\
-5 & 0.5 & 0
\end{array}\right| & X=\left|\begin{array}{ccc}
1 & 0 & 1 \\
0 & 1 & 1 \\
1 & 0 & 1 \\
0 & 1 & 1
\end{array}\right|
\end{array}
$$

## Alternative similarity measure for

0-1 entries: co-occurence

- When $X$ has just 0 or 1 for each entry
- Instead of computing actual covariances from W, compute a similarity score based on count of co-occurrence in $X$


## $\left.\Xi=\begin{array}{lll}1 & 0 & 1 \\ 0 & 1 & 1 \\ 1 & 0 & 1 \\ 0 & 1 & 1\end{array} \right\rvert\,$

- Co-occur(It1, It2) $=0$
- Co-occur(It1, It3) = 2


## Generalization of co-occurrence

## similarity: Association Rules

- From a database of purchases, can find significant co-occurence rules, e.g., person who buys bread and butter => 90\% chance of also buying milk
- It's possible to precompute these association rules (Agarwal et al)


## User-User vs. Item-Item

- Compute pairwise correlations between users
- Compute pairwise correlations between


## $X X^{T}$

 items
## Computational Complexity

- With n items, m users,
- user-user algorithm (unoptimized): about $\mathrm{m}^{2} \mathrm{n}$ operations
- item-item algorithm (unoptimized): about $\mathrm{mn}^{2}$ operations
- \#items may be < \#users
- item-item similarities may be stable over long periods of time => batch computing leads to less inaccuracy


## Predicted Scores for Target Item

- User-user
- Weighted average of other user's ratings of this item
- Weights taken from user-user similarities
- Item-item
- Weighted average of this user's ratings of other items
- Weights taken from item-item similarities


## Finding Items from Items

- Item-item algorithm
- Single starting item
- Find other items with highest correlation
- Starting from a group of items
- Union of results for each item
- (Why are association rules better than the itemitem similarity matrix?)
- User-user algorithm
-??


## Finding Users from Users

- User-user algorithm
- Find other users with highest correlation
- Item-item algorithm
-??


## Web search as a recommender

- Use links between pages as implicit "ratings"
- No separate categories of users, items
- can't easily use user-user algorithm, etc.
- How are the "best" pages for a query recommended?


## Model

- Page is a node
- html link defines a directional link in the graph
- Terminology
- If A has an html to B
- A has an outgoing link to $B$
- $B$ has an incoming link from $A$


## PageRank

- Google's big original idea [Brin \&Page, 1998]
- Idea: ranking is based on "random web surfer":
- start from any page at random
- pick a random link from the page, and follow it
- repeat!
- ultimately, this process will converge to a stable distribution over pages (with some tricks...)
- most likely page in this stable distribution is ranked highest
- Strong points:
- Pages linked to by many pages tend to be ranked higher (not always)
- A link ("vote") from a highly-ranked page carries more weight
- Relatively hard to manipulate


## PageRank, examples



Final distribution properties:
(a) Total weight $=100 \%$
(b) Weight of node is divided among outgoing links.
(c) Weight of node is sum of incoming link weights.

## PageRank, examples



Final distribution properties:
(a) Total weight $=100 \%$
(b) Weight of node is divided among outgoing links.
(c) Weight of node is some of incoming links

## PageRank, mathematically

- Let the stable probabilities be $x_{i}$ for page $i, x_{i}>=0$
- For each $i, j$, define $a_{i j}$ as
- If $j$ links to $i, a_{i j}=(1 /$ number of links of $j)$
- If $j$ does not link to $i, a_{i j}=0$
- Form $A=$ square matrix of $a_{i j}$ for all $i, j$.
- Then, the PageRank probabilities satisfy

$$
A x=x
$$

- $x$ is the eigenvector of the link matrix, with eigenvalue 1
* May need to modify $A$ slightly to ensure unique solution


## Finding the PageRank eigenvector

- One approach: solve linear equation $(A-I) x=\left(\begin{array}{llll}0 & 0 & 0 & .\end{array} 00\right)^{\top}$
- Alternative "power method" is more efficient in practice:
- Start with an arbitrary $X$
- Compute $A^{x}, A^{2} x, \ldots A^{t x}$ (t large)
- $\boldsymbol{A}^{t} \boldsymbol{x}$ is approximately proportional to the correct solution!


## Aside: why the power method works (optional)

- Known: the link matrix A has
- eigenvalue 1 for the correct eigenvector $v^{*}$
- all other eigenvalues $\lambda$ have $|\lambda|<1$
- Known: any $\boldsymbol{x}$ can be expressed as a sum of eigenvectors of $A$
$\mathbf{x}=\mathrm{a}_{0} \mathbf{v}^{*}+\mathrm{a}_{1} \mathbf{v}_{1}+\mathrm{a}_{2} \mathbf{v}_{2}+.$.
- Multiplying by $\mathrm{A} t$ times,
$A^{t x}=a_{0} \mathbf{v}^{*}+a_{1}\left(\lambda_{1}\right)^{t} \mathbf{v}_{1}+a_{2}\left(\lambda_{2}\right)^{t} \mathbf{v}_{2}+.$.
but $\left(\lambda_{1}\right)^{t}$ etc. are very close to 0 for large $t$


## A Sample Graph



$$
A=\begin{array}{cccc}
0 & 0 & .5 & 0 \\
.5 & 0 & 0 & 0 \\
.5 & 1 & 0 & 0 \\
0 & 0 & .5 & 0
\end{array}
$$

## Handling Loops

- Let E be a set of "source" weight ranks
- At each node, random surfer goes to nodes with probabilities in E
- Each node's final rank is a scaled multiple of
- It's source rank PLUS
- The sum of the rank on its backlinks
- Scale it such that the sum of final ranks is 1


## A Sample Graph



$$
E=\begin{aligned}
& .1 \\
& .1 \\
& .1 \\
& .1
\end{aligned}
$$

## Some Intuitions



- Will D's Rank be more or less than $1 / 4$ ?
- Will C's Rank be more or less than B's?
- How will A's Rank compare to D's?


## Mathematical Expression

$$
R^{\prime}=c(A R+E) \quad\|P\|=\left|P_{l}\right|=1
$$

## Power Method Algorithm

Multiply by A, and then normalize so that the sum is 1

$$
R_{i+1}=\frac{A R_{i}+E}{\left|A R_{i}+E\right|}
$$

## Before the First Iteration

## - S

```
r1 . 3
r2 .1
r3 . }
r4 .1
```


## First Iteration

- AR+E
r1. 25
r2 . 25
r3 . 35
r4. 25
- Normalize so sum is 1 (divide by 1.1)

$$
\begin{array}{ll}
\text { r1 } & .22727273 \\
\text { r2 } & .22727273 \\
\text { r3 } & .31818182 \\
\text { r4 } & .22727273
\end{array}
$$

## Second Iteration

## - AR+E

r1 . 25909091
r2 . 21363636
r3 . 44090909
r4. 25909091

- Normalized (divide by 1.17)
r1. 22093023
r2 . 18217054
r3 . 37596899
r4. 22093023


## Third Iteration

- AR+E
r1 . 2879845
r2 . 21046512
r3 . 39263566
r4 . 2879845
- Normalized (divide by 1.18)
r1 . 24424721
r2 . 17850099
r3 . 3330046
r4 . 24424721


## What If More Weight in E?

- Try (1 11 1) instead of (. . . . . .1 1)

| r1 | .23825503 |
| :--- | :--- |
| r2 | .2360179 |
| r3 | .28747204 |
| r4 | .23825503 |

- Try (10 1010 10)
r1 . 24848512
r2 . 24845498
r3 . 25457478
r4. 24848512


## Personalized PageRank

- Pick E to be some sites that I like
- My bookmarks
- Links from my home page
- Rank flows more from these initial links than from other pages
- But much of it may still flow to the popular sites, and from them to others that are not part of my initial set

